

Higher-Order Euclidean Sets

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Example 1: E(7, 12)

Clough and Douthett 1991

Method for computing a maximally even set

$$E_\alpha(k_1, n_1) = \left\{ \left\lfloor \frac{in_1}{k_1} \right\rfloor : i = 0, 1, \dots, k_1 - 1 \right\}$$

n	k	i	calculate	round down
12	7	0	0	0
12	7	1	1.71428571	1
12	7	2	3.42857143	3
12	7	3	5.14285714	5
12	7	4	6.85714286	6
12	7	5	8.57142857	8
12	7	6	10.2857143	10

Toussaint 2005

Geometric method for computing a Euclidean set

Step 1: x x x x x x x

Step 2: x x x x x x x

Step 3: x x x x x

 x x

Step 4: x x x
 . . .
 x x
 x x
 . .

Read down: [0 2 3 5 7 8 10]

-2 ↓

Transpose: [0 1 3 5 6 8 10]

Example 2: E(7,12): Ableton Live

Top part:
[0 1 3 5 6 8 10]

Bottom part:
[0 2 3 5 7 8 10]

These two cycles are the same underlying structure, displaced by 2 time points.



The screenshot shows the Ableton Live interface for a MIDI clip named "1 Claves A 2". The clip is set to a 3/4 time signature and a key signature of one sharp (F#). The pattern is a Euclidean rhythm with 12 steps, 7 notes, and a 1/16 division. The notes are D1 and G0. The pattern is displayed in a piano roll view, showing the notes on a grid. The top part of the pattern is highlighted in yellow, and the bottom part is highlighted in blue. The interface also shows the transport controls, the mixer, and the clip's properties.

Transport: Link, Tap, 90.00, 3 / 4, 0%, 1 Bar, F# C Major, 2. 1. 1

Clip: 1 Claves A 2, All Ins: 1, All Channels: 0, In: Auto, Off, Main: -22.5

Clip Properties: Start: 1. 1. 1, End: 1. 4. 1, Duplicate, Loop, Position: 1. 1. 1, Length: 0. 3. 0, Signature: 4 / 4, Groove: Swing 16ths, Scale: F# C Major

Pattern: Euclidean, Rotation: 0, 2, 0, 0, Steps: 12, Density: 7, Division: 1/16, Auto, Generate

Piano Roll: D1, G0, Velocity: Randomize 100, Ramp: 100, 127, Deviation: 0

Example 4: Clave – rotate first algorithm

$$266: \text{HE} (16 14^0 5^9) = [0 3 6 10 12] < 3 3 4 2 4 >$$

$$394: \text{HE} (16 14^8 5^9) = [0 3 6 10 12] < 3 3 4 2 4 >$$

$$569: \text{HE} (16 13^3 5^8) = [0 3 6 10 12] < 3 3 4 2 4 >$$

$$649: \text{HE} (16 13^8 5^8) = [0 3 6 10 12] < 3 3 4 2 4 >$$

$$676: \text{HE} (16 13^{10} 5^3) = [0 3 6 10 12] < 3 3 4 2 4 >$$

$$756: \text{HE} (16 13^{15} 5^3) = [0 3 6 10 12] < 3 3 4 2 4 >$$

$$801: \text{HE} (16 12^2 5^0) = [0 3 6 10 12] < 3 3 4 2 4 >$$

$$813: \text{HE} (16 12^2 5^{12}) = [0 3 6 10 12] < 3 3 4 2 4 >$$

$$865: \text{HE} (16 12^6 5^0) = [0 3 6 10 12] < 3 3 4 2 4 >$$

$$877: \text{HE} (16 12^6 5^{12}) = [0 3 6 10 12] < 3 3 4 2 4 >$$

$$929: \text{HE} (16 12^{10} 5^0) = [0 3 6 10 12] < 3 3 4 2 4 >$$

$$941: \text{HE} (16 12^{10} 5^{12}) = [0 3 6 10 12] < 3 3 4 2 4 >$$

$$993: \text{HE} (16 12^{14} 5^0) = [0 3 6 10 12] < 3 3 4 2 4 >$$

$$1005: \text{HE} (16 12^{14} 5^{12}) = [0 3 6 10 12] < 3 3 4 2 4 >$$

Example 6: Clave – subtract first algorithm

$$363: \text{HE} (16 14^6 5^{10}) = [0 3 6 10 12] < 3 3 4 2 4 >$$

$$491: \text{HE} (16 14^{14} 5^{10}) = [0 3 6 10 12] < 3 3 4 2 4 >$$

$$628: \text{HE} (16 13^7 5^3) = [0 3 6 10 12] < 3 3 4 2 4 >$$

$$667: \text{HE} (16 13^9 5^{10}) = [0 3 6 10 12] < 3 3 4 2 4 >$$

$$708: \text{HE} (16 13^{12} 5^3) = [0 3 6 10 12] < 3 3 4 2 4 >$$

$$747: \text{HE} (16 13^{14} 5^{10}) = [0 3 6 10 12] < 3 3 4 2 4 >$$

$$801: \text{HE} (16 12^2 5^0) = [0 3 6 10 12] < 3 3 4 2 4 >$$

$$865: \text{HE} (16 12^6 5^0) = [0 3 6 10 12] < 3 3 4 2 4 >$$

$$929: \text{HE} (16 12^{10} 5^0) = [0 3 6 10 12] < 3 3 4 2 4 >$$

$$993: \text{HE} (16 12^{14} 5^0) = [0 3 6 10 12] < 3 3 4 2 4 >$$

Example 7. A “Slate” for clave timeline calculation of depth 2 ($p = 2$)

	n	k_1	k_2	<i>number of sets</i>
Row 0:	16	15	5	$16 \times 16 = 256$
Row 1:	16	14	5	256
Row 2:	16	13	5	256
Row 3:	16	12	5	256
Row 4:	16	11	5	256
Row 5:	16	10	5	256
Row 6:	16	9	5	256
Row 7:	16	8	5	256
Row 8:	16	7	5	256
Row 9:	16	6	5	256

====

sum: 2560

Unfortunately, complexity increases at $O(n^p)$.

Example 9. Deriving a fifth-order Euclidean set from the Clapping Music timeline

Hyper Euclidean Set Generator Subtract first method

Global Controls
 tempo: 120 span (n): 12
 clock numbers on/off:
 pre-sets:
 > 1 Clap Music (a)

Order	Onsets (k)	Rotate (r)	Note	Onsets	IOI	Stdev
1	11	2	C4	0 2 3 4 5 6 7 8 9 10 11	2 1 1 1 1 1 1 1 1 1 1	0.28748
2	8	7	G4	0 1 2 4 5 7 9 10	1 1 2 1 2 2 1 2	0.5
3	7	0	A#4	0 1 2 4 5 7 9	1 1 2 1 2 2 3	0.699854
4	6	0	D5	0 1 2 4 5 7	1 1 2 1 2 5	1.414214
5	3	0	F#5	0 2 5	2 3 7	2.160247

Diagram description: The interface shows five rows of circular diagrams. Each row represents a different order of a Euclidean set. The left diagram in each row shows the onsets (k) on a circle, and the right diagram shows the rotated set (r). The number of onsets (k) decreases from 11 to 3, and the rotation (r) values are 2, 7, 0, 0, 0. To the right of each diagram is a musical staff showing the corresponding rhythm pattern. The tempo is set to 120, and the span (n) is 12. The note for each row is C4, G4, A#4, D5, and F#5 respectively. The IOI (Inter-Onset Interval) and Stdev (Standard Deviation) values are also displayed for each row.

Example 10. "New Clapping Music"

New Clapping Music

Paul V. Miller (2024)

Part 1: 12/8, whole rest.

Part 2: 12/8, whole rest, then rhythmic pattern starting at measure 4. Dynamics: *mf*.

Part 3: 12/8, rhythmic pattern starting at measure 1. Dynamics: *mf*.

Part 4: 12/8, rhythmic pattern starting at measure 1. Dynamics: *mf*.

Measure 4: *f* dynamic.

Measure 7: *mf* dynamic.

Measure 10: *mf* dynamic.

Measure 13: *f* dynamic.

Measure 16: *mf* dynamic.

Measure 19: *mf* dynamic.

Measure 10: *f* dynamic. A box labeled "Repeat 4x" highlights a specific rhythmic pattern.

Measure 13: *f* dynamic.

Measure 16: *mf* dynamic.

Measure 19: *mf* dynamic.

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Appendix 1: Producing higher-order sets

The maximally-even algorithm (Clough and Douthett 1991) produces a conventional first-order Euclidean set E_α where n is the set's span, and k represents the number of desired onsets:

$$E_\alpha(k_1, n_1) = \left\{ \left\lfloor \frac{in_1}{k_1} \right\rfloor : i = 0, 1, \dots, k_1 - 1 \right\}$$

E_α is a set of onsets where j is the index of each e :

$$e_{\alpha_j} \in E_\alpha : j = 1, 2, \dots, k_1$$

To produce a **second-order Euclidean set** E_2 , let $n_2 = k_1$, and k_2 equal the number of desired onsets in the second iteration. Then by simple recursion,

$$E_\beta(k_2, n_2) = \left\{ \left\lfloor \frac{in_2}{k_2} \right\rfloor : i = 0, 1, \dots, k_2 - 1 \right\}$$

produces an **index set** such that each $e_{\beta_j} \in E_\beta$ is the j th term in E_α . The set of e_β onsets in E_α is the second-order Euclidean set E_2 . Consequently, E_2 is *not* necessarily a subset of E_α .

Appendix 3. The harmonic minor scale as a second-order Euclidean set

Hyper Euclidean Set Generator

Subtract first method

Global Controls

tempo

▶ 120

span (n)

▶ 12

clock numbers on/off

pre-sets
:

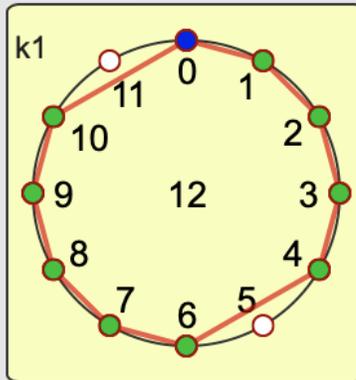


▶ 14

harm min scale

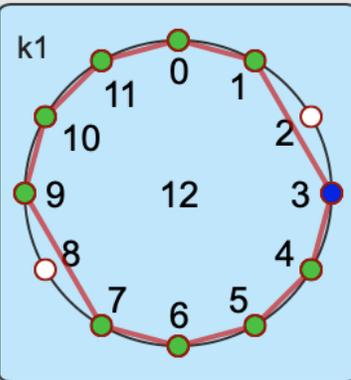
onsets (k1)

▶ 10



rotate (r1)

▶ 3



note ▶ C4

onsets 0 1 3 4 5 6 7 9 10 11

IOI 1 2 1 1 1 1 2 1 1 1

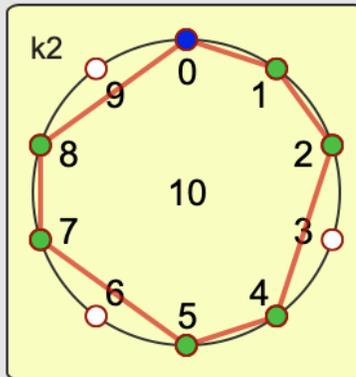
play stop

stdev 0.4



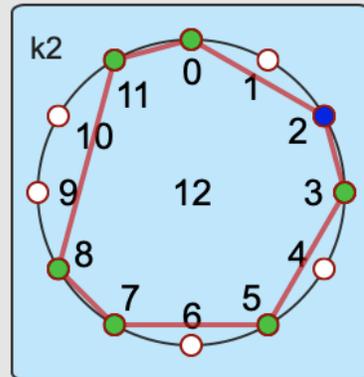
onsets (k2)

▶ 7



rotate (r2)

▶ 2



note ▶ G4

onsets 0 2 3 5 7 8 11

IOI 2 1 2 2 1 3 1

play stop

stdev 0.699854

